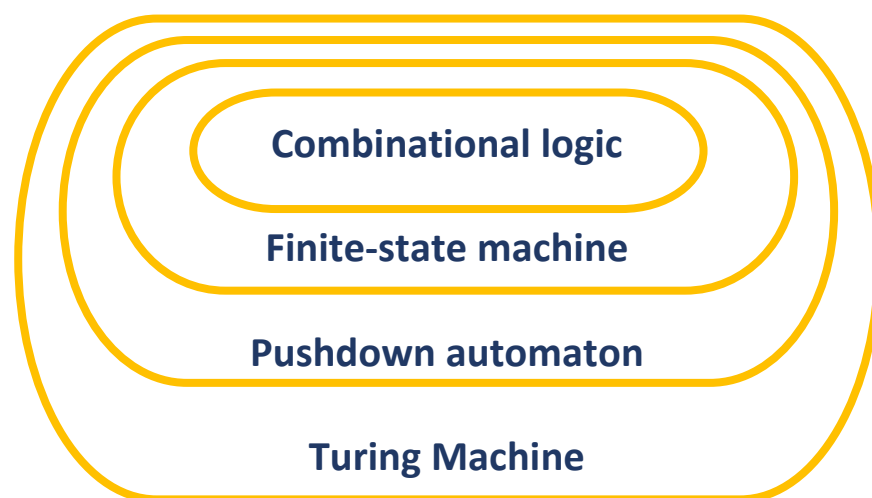


## Turing Machine

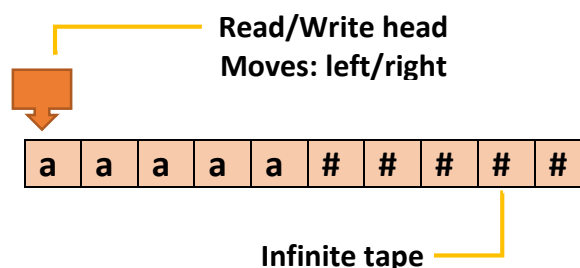
A Turing machine is a mathematical model of computation describing an abstract machine that manipulates symbols on a strip of tape according to a table of rules. Despite the model's simplicity, it is capable of implementing any computer algorithm.



### Abilities of Turing machine

1. A Turing machine is similar to a DFA or PDF
2. It has the following abilities:
  - It can **read** or **write** to tape.
  - It can **move left** or **move right** on the tape.
  - It **halts** as soon as it reaches either the special **accept state** or the special **reject state**

### Scheme of the Turing machine

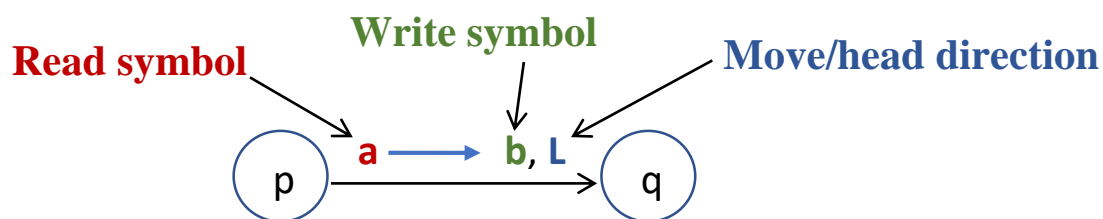


### The formal definition of TM

A deterministic Turing Machine is a tuple consisting of several objects:

- $Q$  is a finite set called the states.
- $\Sigma$  is a finite alphabet, a finite set not containing the **blank** symbol
- $\Gamma$  is the tape alphabet, where  $\Sigma \subseteq \Gamma$  and  $\epsilon \in \Gamma$
- $\delta: Q \times \Gamma \rightarrow \delta(Q \times (\Gamma \times \{L,R\}))$  is a function called the transition function
- $S$  is an element of  $Q$  called the start state
- $F$  is a subset of  $Q$  called the accepted states
- TM  $M = (Q, \Gamma, b, \Sigma, \delta, S, F)$

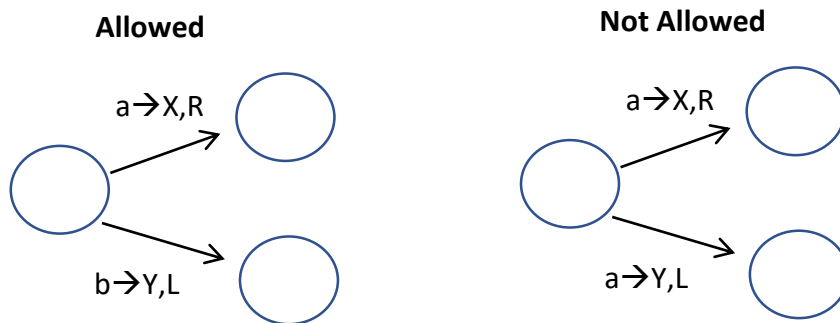
### Transitions



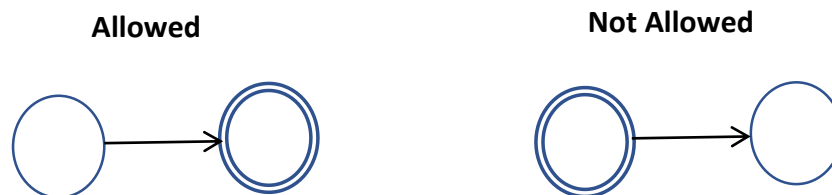
$$\delta(p, a) = (q, b, L)$$

## The Language TM Accepts

- A Turing Machine accepts its input if it reaches an accepting configuration.
- The set of inputs it accepts is called its **language**.
- Turing Machine are deterministic



- Final states



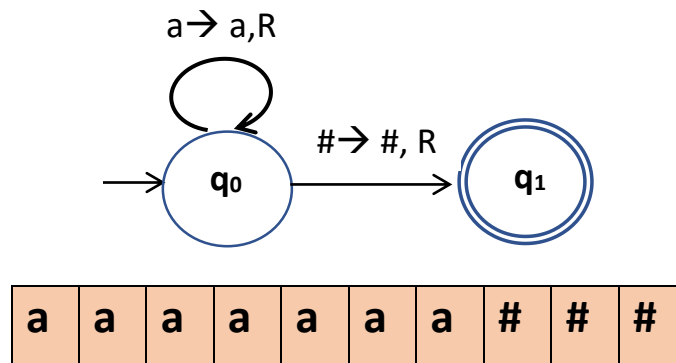
- Final states have no outgoing transitions.
- In a final state the machine halts.

- Acceptance

- **Accept input:** If machine halts in a final state.
  - **Reject input:** If machine halts in a non-final state
- Or**
- If machine enters an infinite loop.

**TM example1:**

Turing machine that accepts language  $a^*$

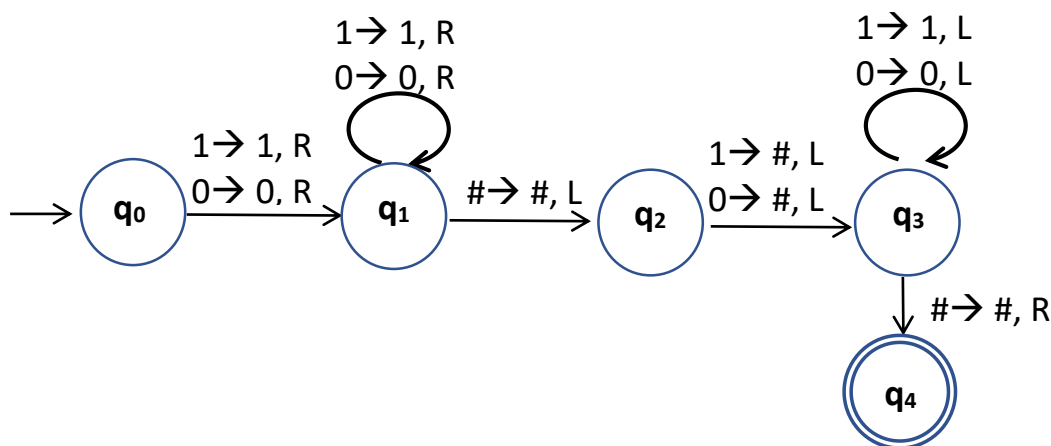
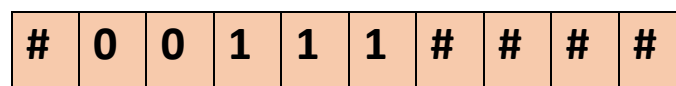


Ex2: Turing Machine to erase the first symbol in the right side

Sol:

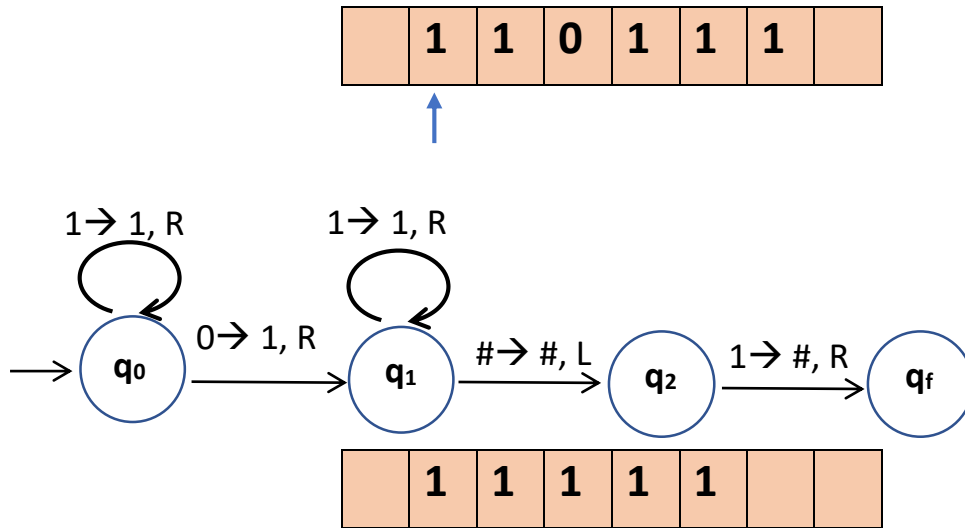
**Steps**

1. Move the head in the right direction
2. Read symbols without change
3. When reading the symbol # move to the left
4. Change the symbol into # and move to the left
5. Move the head in the right direction and read without change
6. When reach the first symbol in the left move to the right and stop.

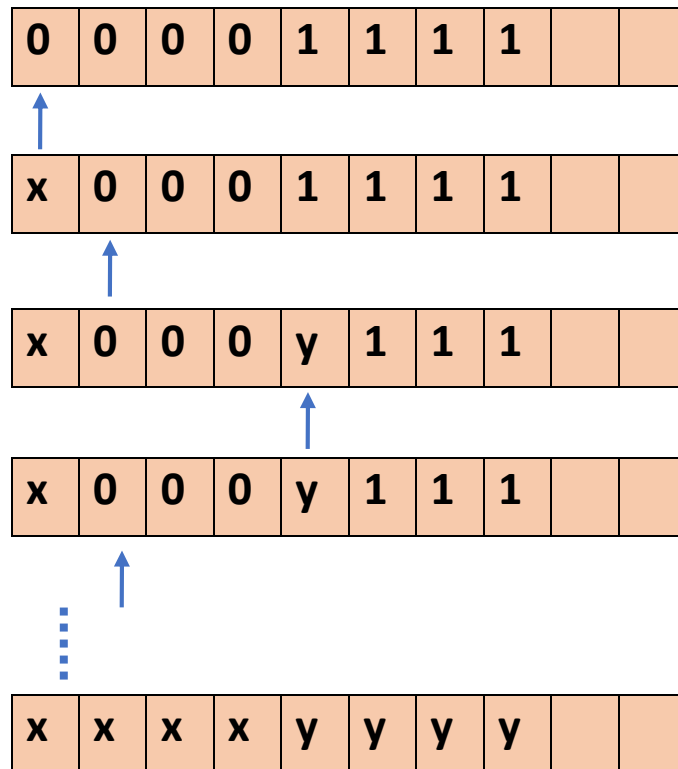


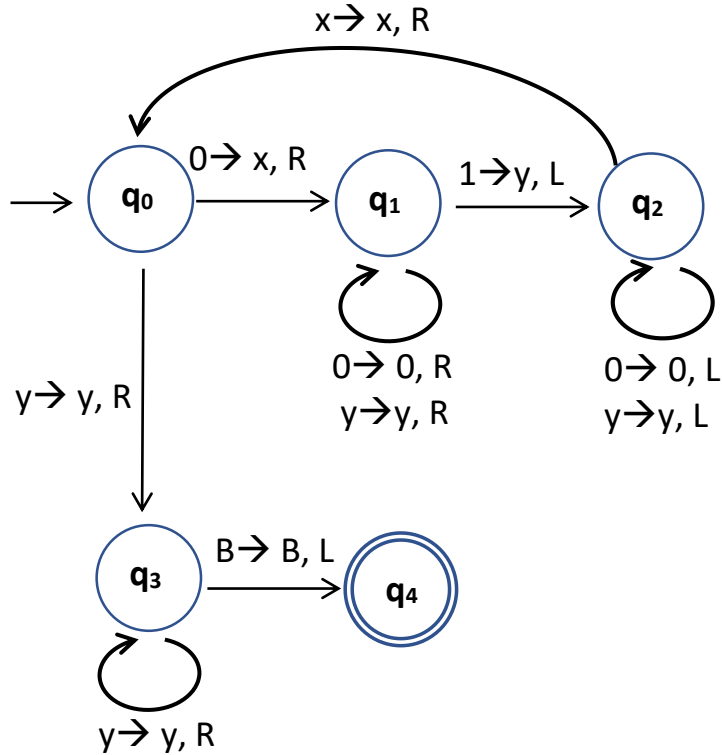
Ex2: design TM for binary addition is computable

Sol: test input (2,3) = 110111



Ex3: A Turing Machine M that accepts the language  $L = \{0^n 1^n | n \geq 0\}$





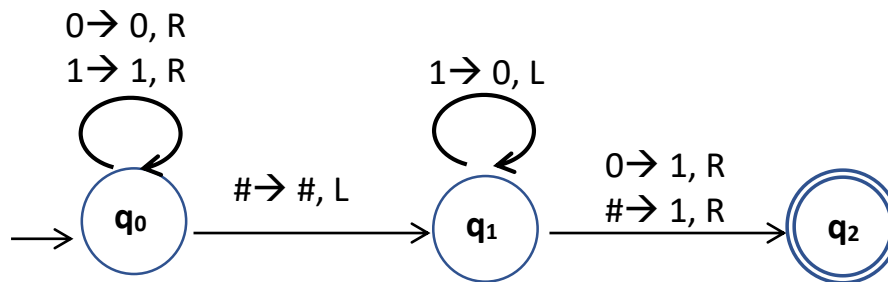
Ex3:

Design a TM to add 1 to a binary number  $x$ .

carry	0	1	1	0
X	0	0	1	1
				1
result	0	1	0	0

#### خطوات الحل

1. نقرأ الآلة رمز دون تغييره وتتحرك نحو اليمين.
2. تكرر الخطوة 1 حتى تصل الى الرمز #.
3. تتحرك خطوة واحدة نحو اليسار تقف عند اول مرتبة من جهة اليمين least significant bit
4. اذا كان الرمز 0 غيره الى 1 وتوقف .
5. اذا كان الرمز 1 غيره الى 0 وانتقل خطوة نحو اليسار.
6. كرر الخطوات 4 و5 حتى تصل الى الرمز # الذي يعني انتهاء العملية والتوقف عندها.



➤  $X = 0110$

#	0	1	1	0	#	#	#	#	#
#	0	1	1	1	#	#	#	#	#

➤  $X = 111$

#	0	1	1	1	#	#	#	#	#
#	1	0	0	0	#	#	#	#	#