

## The generalized $k$ -connectivity of equally complete bipartite graphs and their line graphs

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### Abstract

Generalized  $k$  connectivity for graphs and generic graphs is recognized as NP-complete, a parameter that measures the network's ability to connect vertices. Suppose  $k_{m,n}$  a completely connected bipartite graph represents the maximum number of internally disjoint Steiner trees (IDSTs) joining a subset  $S \subseteq V(G)$  of  $k$  vertices in  $G$ . In this context, Steiner trees (or "S -trees")  $T_1, T_2$  are considered internally detached if and only if  $V(T_1) \cap V(T_2) = S$  and  $E(T_1) \cap E(T_2) = \phi$ . We define generalized  $k$ -connectivity as  $\kappa_k(G)$ . This study focuses on calculating the precise values of generalized  $k$ -connectivity for line- graph of bipartite graphs with  $k = 3, 4$  and generalized  $k$ -connectivity for bipartite graphs with  $k \geq 3$ .

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### 1. Introduction

A desired pair  $(V(G), E(G))$  represents  $(G)$  graph, where both of  $(V(G), E(G))$  symbolize vertices, edges sets, respectively. By the number of vertices which is directly adjacent to  $x \in V(G)$  degree of a vertex is determined it that symbolize  $d$  by  $d_G(x)$ . The minimum degrees of  $G$  are represented by  $\delta = \delta(G)$ , respectively. When  $\delta$  is equal to some natural number  $q$ , the graph is

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called  $q$  -regular. When the vertices of  $(G)$  are separated into the sub-sets  $X, Y$ , then each edge has an end vertex at  $X$  and another at  $Y$  [2]. A completed bipartite -graph, symbolized as  $(k_{m,n})$  that is a bipartite graph where all vertices in  $X$  are adjacent to all vertices in  $Y$ , with  $m = |X|$  and  $n = |Y|$ . The application of graphs see [17,18],  $L(G)$  is the symbol of the line graph of  $G$  where, exactly when the corresponding edges are adjacent that the two vertices in  $L(G)$  are adjacent. The vertices of  $L(G)$  correspond to the edges of  $G$ . Connection is a basic understood in graph theory, measuring the ability  $G$  to remain connected after removing vertices or edges. Connectivity for the graph  $G$  symbols as  $\kappa(G)$ , representing the mini number of vertices which have to remove to separate the graph or reduce it to a trivial state. The  $G_{\kappa} - C$ , introduced by Chartrand et al. in [3], is a stronger measure of connectivity used to assess the dependability and fault tolerance in the network. It considers any subset  $(S)$  of the vertex set  $V(G)$  and defines a Steiner tree as a tree in  $G$  that contains all vertices in  $S$ .  $S$  - trees are internally edge-disjoint trees if they share no common edges. The  $G_{\kappa} - C$  for  $G$ , symbolize as  $\kappa_k(G)$ , is the least connectivity value amongst all subsets  $S$  of  $V(G)$  with  $k$  vertices. J.M. Xu [12] established in 2001 that a graph has at least two satisfying edges. Ch.F et al. [10] found a solution for generalized connectivity in graphs for any two integers  $n$  and  $k$ . The upper and lower bounds of generalized connectivity in a general graph  $G$  are discussed in [4,6]. Sh. Li et al.[7] introduced the  $G_{\kappa} - C$  of star graphs for  $k = 3$ , and later Dh. D. Kadhim and A.A. Najim [2] improved it for an equally complete  $k$  - partite graph in 2019. The generalized 4-connectivity of hypercube for  $CBPG$  was studied by Sh. Lin and Q. Zh [9] . In 2022, S. Li, Zhao, et. al. [15] studied internally disjoint trees in the line graph and total graph of  $CBPG$ . For further information on generalized connectivity, refer to [1,5,8,11,13-16].

**Abbreviations**

Abbreviations	Meaning
$G_{\kappa} - C$	Generalized $\kappa$ -connectivity
IDSTs	internally disjoint $S$ - trees
MNIDST	maximum number of IDSTs
$CBPG$	complete bipartite graph

**2. Preliminary Results**

**Lemma 2.1:** [6] *Any connected  $G$  graph have order  $n$  arranged by minimum degree  $\delta$ , if there are two adjacent vertices with same  $\delta$ , the system have  $k_k(G) \leq \delta - 1$  of  $3 \leq k \leq n$ .*

**Lemma 2.2:** [6] *If  $n$  and  $k$  are integers where  $(2 \leq k \leq n)$ ,  $G$  is connected with  $n$  vertices so, vertex connectivity  $(\kappa_k(G))$  is less or equal than edge connectivity  $(\lambda_k(G))$ , which is in turn less or equal than the less degree of  $(\delta(G))$ .*

**Lemma 2.3:** [8] *For any edge  $e = xy \in E(G)$ , the grade of  $e$  in the line graph  $L(G)$  is equal to the sum of the degrees of  $(x, y)$  in simple undirected graph,*

minus 2. Specifically, whether  $G$  is  $q$ -regular, so  $L(G)$  is  $(2q - 2)$ -regular. In fact, the IDSTs with the maximum number of connections, denoted as MNIDST for convenience, are of interest. The count of cases for an equally CBPG can be calculated as the ceiling of  $(k + 1)/2$ , whereas for a non-equally CBPG, it is simply  $k + 1$ . Furthermore, any vertex with a non-positive index is not part of the vertex set  $V(G)$ , and consequently, any edge incident to such a vertex is not included in  $E(G)$ .

### 3. Principal findings

This section's goal is to assess the  $G_\kappa - C$  of both the equally CBPG, considering a range of values for  $k$  from 3 to 9. Moreover, we will examine the non-equally CBPG along with its line, focusing specifically on the cases when  $k$  takes on the values 3, 4, and 5.

**Theorem 3.1:** Any integer  $n$  and  $k$ , where  $n \geq 2$  and  $k$  takes the values 3, 4, 5, the following statement holds true: the rank of CBPG  $k_{n,n}$  is equal to  $n - 1$ .

**Proof:** Consider the CBPG  $k_{n,n}$  with vertex sets  $(X = \{x_1, x_2, \dots, x_n\}, Y = \{y_1, y_2, \dots, y_n\})$  where each vertex in  $X$  is adjacent to every vertex in  $Y$ , and  $|X| = |Y| = n$ . Since  $k_{n,n}$  is a regular graph, the minimum degree is  $\delta = n$ . By applying Lemma (2.1), we obtain the inequality.  $\kappa_k(k_{n,n}) \geq n - 1$ . To establish the lower bound  $\kappa_k(k_{n,n}) \geq n - 1$ , we require to demonstrate for every  $k$ ,  $S \subseteq V(k_{n,n})$  at least there is  $n - 1$  IDSTs connecting the vertices in  $S$  within  $k_{n,n}$ , where  $|S| = k$ . we discuss the case define as:

**Case(r)**  $S_r = \{y_{r-3}, y_{r-2}, \dots, y_{r-1}, x_1, x_2, \dots, x_{k-r+1}\}$ , where  $r = 1, 2, 3$ , then MNIDST in  $k_{n,n}$  is  $n - 1 \leq \kappa_k(S) \leq n$ , as explained in:  $T_i = x_1 y_i x_2 \cup x_3 y_i x_4 \cup \dots \cup x_{q-1} y_i x_h \cup x_j y_1 x_h y_{3-m}$  (unless  $k = 5, r = 3$ ) and  $i = 3$ ) and  $T_4 = y_1 x_1 y_3 x_2 \cup y_3 x_3 y_2$  at  $(k = 5, r = 3$  and  $i = 3$  Where

$$i = \begin{cases} 1, 2, \dots, n & \text{if } r = 1 \\ 2, \dots, n & \text{if } r = 2, 3 \end{cases}, q = \begin{cases} k - r + 1 & \text{if } r = 1, 2 \\ k - 3 & \text{if } r = 3 \end{cases}, h = \begin{cases} 0 & \text{if } r = 1 \\ 1 & \text{if } r = 2, 3 \end{cases}$$

$$m = \begin{cases} 3 & \text{if } r = 2 \\ 1 & \text{if } r = 3 \end{cases} \text{ and } j = \begin{cases} 0 & \text{if } k = 3, 4, 5, r = 2 \text{ and } i = 2, 3, \dots, n \text{ and} \\ k & \text{if } k = 5, r = 3 \text{ and } \forall i \geq 4 \\ 3 & \text{if } k = 5, r = 3 \text{ and } \forall i = 2 \end{cases}$$

When  $h = 0$  then the branch  $y_i x_h$  is not exist. From the argument above, we get  $n - 1 \leq \kappa_k(S) \leq n$ . Thus, we get  $\kappa_k(k_{n,n}) = n - 1$ .

**Lemma 3.1:** Let  $G = k_{n,n}$  where  $n \geq 3$ , we have  $\kappa_k(k_{n,n}) \leq \delta - 2$  where  $k \geq 6$ .

**Proof:** Let  $k_{n,n}$  be a CBPG,  $G(X \cup Y, E)$ , such that  $|X| = |Y| = n$ . Let  $S \subseteq V(k_{n,n})$  where  $\delta = k$ ,  $S = \{u_1, u_2, \dots, u_k\}$  where  $u_1, u_2 \in Y$  and  $u_i \in X$ ;  $i = 3, 4, \dots, k$ . We want to prove that the MNIDSTS connecting in  $G$  are  $\delta - 2$ , which is  $T_1, T_2, \dots, T_{\delta-2}$ . Clearly, the Steiner tree containing the set  $S$  has a

maximum degree of two or more. Assuming that  $d_G(u_1) \geq 3$  in one particular Steiner tree, denoted as  $T_1$ , implies the existence of  $\delta - 3$  Steiner trees connecting  $S$  in the graph  $G$ , distinct from  $T_1$ . Consequently, the minimum number of IDSTs (MNIDSTs) connecting  $S$  in  $G$ . Alternatively, if we suppose that  $d_G(u_1) = 2$  in two Steiner trees, say  $T_1$  and  $T_2$ , then there exist  $\delta - 4$  additional Steiner trees connecting  $S$ , resulting in  $\delta - 2$  MNIDSTs connecting  $S$  in  $G$ . In the case where the maximum degree of vertices in  $S$  is two or less in a single tree, denoted as  $T_1$ , and the vertices in  $S$  in other trees have a degree of one,  $T_1$  requires at least  $k - 5$  vertices from the set  $Y$  (excluding  $u_1$  and  $u_2$ ). Therefore, the remaining number of vertices from the set  $Y$ , which are different from  $u_1, u_2$ , and the  $k - 5$  vertices, is  $n - k - 3$ . In this scenario, the MNIDSTs connecting  $S$  in  $G$  amount to  $n - k - 3$ . Based on the above results, we can conclude that  $\kappa_k(k_{n,n}) \leq \delta - 2$ .

**Theorem 3.2:** Any integer  $n, k$  where  $n \geq 3$  and  $k = 6, 7, 8, 9$ . Then  $\kappa_k(k_{n,n}) = n - 2$ .

**Proof:** Let  $k_{n,n}$  be a CBPG,  $G(X \cup Y, E)$  in which each vertex of  $X$  is adjacent with each vertex of  $Y = \{y_1, y_2, \dots, y_n\}$  such that  $|X| = |Y| = n$ . Since  $k_{n,n}$  is an regular graph this means that  $\delta = n$  then by lemma (3.1) we get  $\kappa_k(k_{n,n}) \leq n - 2$ . To prove  $\kappa_k(k_{n,n}) \geq n - 2$ , the suggested system require that for any  $k$ -subset  $S \subseteq V(k_{n,n})$  at least there is an  $n - 2$  IDSTs connecting  $S$  in  $k_{n,n}$ . we discuss five cases as define:

**Case(r)**  $S_r = \{y_{r-3}, y_{r-2}, y_{r-1}, x_1, x_2, \dots, x_{k-r+1}\}$ , where  $r = 1, 2, 3$ , then the MNIDST connecting Sink  $k_{n,n}$  is  $n - 2 \leq S \leq n$ , as explained in:

$$T_j = x_1 y_i x_2 \cup x_3 y_i x_4 \cup \dots \cup x_q y_i x_h \cup y_1 x_h y_{3-m}$$

$$\text{Where } i = \begin{cases} 1, 2, \dots, n & \text{if } r = 1 \\ 2, 3, \dots, n & \text{if } r = 2 \\ 3, 4, \dots, n & \text{if } r = 3 \end{cases} \quad j = \begin{cases} i & \text{if } r = 1 \\ i - 1 & \text{if } r = 2, q = \begin{cases} k & \text{if } r = 1 \\ k - r & \text{if } r = 2, 3 \end{cases} \\ i - 2 & \text{if } r = 3 \end{cases}$$

$$h = \begin{cases} 0 & \text{if } r = 1 \\ i & \text{if } r = 2, 3 \end{cases} \quad \text{and } m = \begin{cases} 3 & \text{if } r = 2 \\ 1 & \text{if } r = 3 \end{cases}$$

When  $h = 0$  or negative the branch  $y_i x_h$  is not exist.

**Case(4)**  $S_4 = \{y_1, y_2, y_2, x_1, x_2, \dots, x_{k-3}\}$ . Then the MNIDST in  $k_{n,n}$  is  $n - 2$ , as explained in :  $T_1 = y_2 x_1 y_1 x_2 y_3 x_{m+2} \cup y_1 x_m \cup x_3 y_2 x_{m+1}, T_2 = x_1 y_4 x_2 y_2 \cup x_3 y_4 x_4 \cup \dots \cup x_{k-4} y_4 x_{k-3} \cup y_1 x_3 y_2$ , where  $x_m, x_{m+1}, x_{m+2} \in S$  and

$$m = \begin{cases} 0 & \text{if } k = 6 \\ 4 & \text{if } k = 7, 8, 9 \end{cases}, T_3 = x_1 y_5 x_2 \cup x_3 y_5 x_4 y_2 \cup y_3 x_5 y_5 x_6 y_1$$

(used only at  $k = 8, 9$ ),  $T_4 = x_1 y_6 x_2 \cup x_3 y_6 x_4 y_3 \cup y_1 x_5 y_6 x_6 y_2$  (used only at  $k = 9$ ) and  $T_{i-2} = x_1 y_i x_2 \cup x_3 y_i x_4 \cup \dots \cup x_{k-3} y_i x_i y_1 \cup y_2 x_i y_3$  Where

$$i = \begin{cases} k - 2, k - 1, \dots, n & \text{if } k = 7, 8, 9 \\ k - 1, k, \dots, n & \text{if } k = 6 \end{cases}$$

**Case (5)**  $S_4 = \{y_1, y_2, y_2, y_4, x_1, x_2, \dots, x_{k-4}\}$ , (used only at  $k = 8, 9$ ). Then the *MNIDST* in  $k_{n,n}$  is  $n - 2$ , as explained in:  $T_1 = x_4y_1x_1y_3x_2x_m \cup y_3x_3y_4$ ,  $T_2 = x_3y_2x_1y_4x_2y_1x_m \cup y_4x_4y_3$   $T_3 = x_1y_5x_2 \cup y_1x_3y_5x_4y_2 \cup y_3x_5 \cup x_5y_4$  and  $T_{i-2} = x_1y_i x_2 \cup x_3y_i x_4 \cup x_m y_i x_i y_1 \cup y_2 x_i y_3 \cup y_4 x_i$  where  $i = 6, 7, \dots, n$  and

$$m = \begin{cases} 0 & \text{if } k = 8 \\ 5 & \text{if } k = 9 \end{cases}. \text{ Thus } \kappa_k(k_{n,n}) = n - 2.$$

**Theorem 3.4:** *The extended 4-connectivity for line G graph of CBPG is  $\kappa_4(L(k_{m,n})) = m + n - 3$  where  $1 \leq n \leq m$ .*

**Prove:** Spouse  $X = \{x_1, x_2, \dots, x_m\}$ ,  $Y = \{y_1, y_2, \dots, y_n\}$  be two parts of  $k_{m,n}$  such  $V(K_{m,n}) = V(X \cup Y)$  and let  $L(k_{m,n})$  be the line graph of a CBPG  $k_{m,n}$  with  $1 \leq n \leq m$ . Since  $L(k_{m,n})$  is a  $m + n - 2$ -regular graph, we get by lemma (2.1)  $\kappa_4(L(k_{m,n})) \leq m + n - 3$ . Next we prove that  $\kappa_4(L(k_{m,n})) \geq m + n - 3$  to achieve this we need to proof that there is IDSTs connecting in for subset. We can write  $L(k_{m,n})$  as  $n$  sets and  $m$  vertices in each set  $x_i y_j \in V(L(k_{m,n}))$  where  $(i = 1, 2, \dots, m)$  and  $(j = 1, 2, \dots, n)$  is a vertex in  $L(k_{m,n})$  and its edge in  $k_{m,n}$ , also, the edge in  $L(k_{m,n})$  are  $(x_i y_j, x_i y_{j'})$  and  $(x_i y_j, x_{i'} y_j)$  where  $(i = 1, 2, \dots, m)$ ,  $(j = 1, 2, \dots, n)$ ,  $(i' = 1, 2, \dots, m)$  and  $(j' = 1, 2, \dots, n)$ .

**Case (1):** If  $u = x_1 y_1$ ,  $v = x_2 y_1$ ,  $r = x_3 y_1$  and  $w = x_4 y_1$ . Then the *MNIDST* in  $L(k_{m,n})$  is  $m + n - 3$ , selected as:  $T_1 = wvur$ ,  $T_2 = uwr v$ ,  $T_{i+1} = ux_i y_i x_2 y_i v \cup x_i y_i x_3 y_i r \cup x_1 y_i x_4 y_i w$  where  $(i = 1, 2, \dots, n)$  and  $T_{n+j-3} = ux_j y_1 v \cup x_j y_1 r \cup x_j y_1 w$  where  $(j = 5, 6, \dots, m)$ .

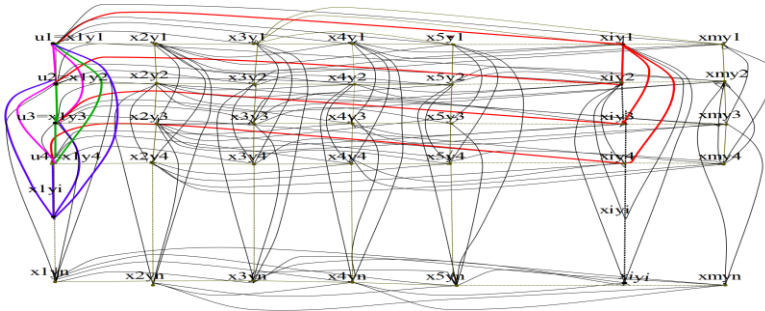
**Case (2):** If  $u = x_1 y_1$ ,  $v = x_2 y_1$ ,  $r = x_3 y_1$  and  $w = x_1 y_2$ . Then the *MNIDST* in  $L(k_{m,n})$  is  $m + n - 3$ , selected as :  $T_1 = wuvr$ ,  $T_2 = vx_2 y_2 w x_3 y_2 r u$ ,  $T_{i-1} = ux_i y_1 v \cup x_i y_1 r \cup x_i y_1 x_i y_2 w$  where  $(i = 4, 5, \dots, m)$  and  $T_{m+j-3} = ux_1 y_j w \cup x_1 y_j x_2 y_j v \cup x_1 y_j x_3 y_j r$  where  $(j = 2, 3, \dots, n)$ .

**Case (3):** If  $u = x_1 y_1$ ,  $r = x_1 y_2$  and  $w = x_2 y_2$ . Then the *MNIDST* in  $L(k_{m,n})$  is  $m + n - 3$ , selected as :  $T_1 = vurw$ ,  $T_{i-2} = ux_i y_1 v \cup x_3 y_1 x_i y_2 r \cup x_i y_2 w$  where  $(i = 3, 4, \dots, m)$  and  $T_{m+j-3} = ux_1 y_j r \cup x_1 y_j x_2 y_j v \cup x_2 y_j w$  where  $(j = 3, 4, \dots, n)$ .

**Case (4):** If  $u = x_1 y_1$ ,  $v = x_2 y_1$ ,  $r = x_1 y_2$  and  $w = x_1 y_3$ . Then the *MNIDST* in  $L(k_{m,n})$  is  $m + n - 3$ , selected as :  $T_1 = rwuv$ ,  $T_2 = ux_2 y_2 v \cup x_2 y_2 x_2 y_3 w$   $T_i = ux_i y_1 v \cup x_i y_1 x_i y_2 r \cup x_i y_1 x_i y_3 w$  where  $(i = 3, 4, \dots, m)$  and  $T_{m+j-3} = ux_1 y_j r \cup x_1 y_j w \cup x_1 y_j x_2 y_j v$  where  $(j = 4, 5, \dots, n)$ .

**Case (5):** If  $u = x_1 y_1$ ,  $v = x_1 y_2$ ,  $r = x_1 y_3$  and  $w = x_1 y_4$ . Then the *MNIDST* in  $L(k_{m,n})$  is  $m + n - 3$ , (see Figure 1) selected as :  $T_1 = wvur$ ,  $T_2 = uwr v$ ,

$T_{i+1} = ux_iy_1x_iy_2v \cup x_iy_1x_iy_3r \cup x_iy_1x_iy_4w$  where  $(i = 2,3, \dots, m)$  and  $T_{m+j-3} = ux_1y_jv \cup x_1y_jr \cup x_1y_jw$  where  $(j = 5,6, \dots, n)$ .



**Figure 1**

**By the five cases above, deduce any sub-set (if IDSTs are connecting). Thus  $\kappa_4(L(k_{m,n})) = m + n - 3$ .**

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