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## Minimum diameter and $t_f(f, q)$

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### Abstract

For given integer  $e$  and  $q$ , let  $p(e, q)$  (res.  $C(e, q)$ ) indicate to (MDe-PD). The maximum diameter deleting  $e$  edges from a connected graph of diameter  $q$  is  $f(e, q)$ , the maximum number  $T_f(f, q)$  of edges that have to be delete from a path  $q$ . In our study show construction  $p(7, 8k - 10) = k$  for  $k \geq 3$  And  $p(7, q') = k$  for  $k \geq 3$  and  $q' = 8k - 6$ , where  $e$  and  $q$  are positive integer also show that  $f(4, q) = 5q - 4$  where  $q \geq 4$  and  $f(7, q) = 8q - 10$  where  $q \geq 5$ . Also

$$1T_f 1(f, q)1 \geq \left\lceil \frac{f-q-3}{q+1} \right\rceil \text{ Where } 1f(e, q)1 \leq 1q(e+1)1 - e + 3 \text{ and}$$

$$\left\lceil \frac{f-6}{q-2} \right\rceil - 1 \geq T_f(f, q) \geq \begin{cases} \left\lceil \frac{f-6}{q+1} \right\rceil & \text{if } f(e, q) \leq (e+1)q + 1 \\ \left\lceil \frac{f-2}{q-2} \right\rceil - 1 & \text{if } f(e, q) \leq (e+1)q - e + 1 \end{cases}$$

$$\text{Where } f(e, q) \geq (e+1)q - 2e + 4$$

**Subject Classification:** 05C10, 05C30, 57M15.

**Keywords:** Length of longest path, Addition and deletion problems.

## 1. Introduction

Let  $V(G)$  is a vertex set and  $E(G)$  is the edges set; A graph  $G = (V, E)$  be an undirected simple graph [2]. Distance between  $w$  and  $z$  in  $G$  is  $d_G(w, z)$

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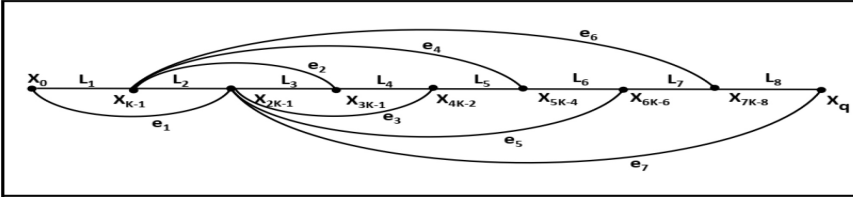
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which is the length of shortest path joining them. Maximum distance between any two vertices of  $G$  is the Diameter which denoted by  $d(G)$  [3]. We can see more in [1,4]. In graph theory :given positive integer  $e$  and  $q$ , let  $F(e, q)$  indicate to the minimum diameter of an altered graph obtained by adding  $e$ - extra edges to a graph with diameter  $q$  [3]. Let  $p(e, q)$  (res.  $C(e, q)$ ) indicate the minimum diameter of a graph obtained by adding  $e$ - edges to a path of length  $q$  (MDe-PD). In [5] we can see  $F(e, q) = p(e, q)$  for any integer  $e$  and  $q$ . The convers to edge-addition problem is called edge deletion problem. Let  $g(e, q)$  denoted the maximum possible diameter of graph  $H$  obtained by deleting  $e$ edges from  $(e + 1)$ - edge-connected graph  $G$  with diameter  $q$

Let  $f(e, q)$  denoted the maximum diameter of a connected graph obtained after deleting  $e$  edges from connected graph with diameter  $q$ . exact value of  $f(e, q)$  have been obtained for some  $e$  or  $q$  [6]. Let  $T_F(f, q)$  be a maximum number of edges that delete from  $q$  ( length of path) to changed it to be a graph of diameter at most  $F$  [7].

**Theorem 2.1 :**  $p(7, 8k - 10) = k$  for  $k \geq 3$

*Proof :* Let  $q = 8k - 10$ . We add seven edges :



Construction of theorem 2.1

$$e_1 = x_0 x_{2k-1}, e_2 = x_{k-1} x_{3k-1}, e_3 = x_{2k-1} x_{4k-2}, e_4 = x_{k-1} x_{5k-4}$$

$$e_5 = x_{2k-1} x_{6k-6}, e_6 = x_{k-1} x_{7k-8}, e_7 = x_{2k-1} x_q$$

To path  $L = x_0 x_1 \dots x_q$  which be a simple path. Seven vertices  $x_{k-1}$ ,  $x_{2k-1}$ ,  $x_{3k-1}$ ,  $x_{4k-2}$ ,  $x_{5k-4}$ ,  $x_{6k-6}$  and  $x_{7k-8}$ , partition  $L$  into eight segments:

$$L_1 = (x_0, x_{k-1}), L_2 = (x_{k-1}, x_{2k-1}), L_3 = (x_{2k-1}, x_{3k-1}), L_4 = (x_{3k-1}, x_{4k-2}),$$

$$L_5 = (x_{4k-2}, x_{5k-4}), L_6 = (x_{5k-4}, x_{6k-6}), L_7 = (x_{6k-6}, x_{7k-8}), L_8 = (x_{7k-8}, x_q)$$

Now, we will define 28 cycles as follows :

$$\begin{aligned}
 C^1 &= L_1 \cup L_2 + e_1 & C^2 &= L_1 \cup L_3 + e_1 + e_2 \\
 C^3 &= L_1 \cup L_5 + e_3 + e_4 + e_1 & C^4 &= L_1 \cup L_8 + e_1 + e_6 + e_7 \\
 C^5 &= L_2 \cup L_3 + e_2 & C^6 &= L_2 \cup L_5 + e_4 + e_3 \\
 C^7 &= L_2 \cup L_7 + e_6 + e_5 & C^8 &= L_3 \cup L_4 + e_3 \\
 C^9 &= L_3 \cup L_5 + e_2 + e_4 + e_3 & C^{10} &= L_3 \cup L_6 + e_2 + e_4 + e_5 \\
 C^{11} &= L_3 \cup L_7 + e_2 + e_6 + e_5 & C^{12} &= L_3 \cup L_8 + e_2 + e_6 + e_7 \\
 C^{13} &= L_4 \cup L_5 + e_2 + e_4 & C^{14} &= L_4 \cup L_6 + e_3 + e_5 + e_4 + e_2 \\
 C^{15} &= L_4 \cup L_7 + e_3 + e_5 + e_2 + e_6 & C^{16} &= L_4 \cup L_8 + e_3 + e_7 + e_2 + e_6 \\
 C^{17} &= L_2 \cup L_8 + e_6 + e_7 & C^{18} &= L_5 \cup L_6 + e_3 + e_5 \\
 C^{19} &= L_5 \cup L_7 + e_3 + e_5 + e_4 + e_6 & C^{20} &= L_5 \cup L_8 + e_3 + e_7 + e_4 + e_6 \\
 C^{21} &= L_6 \cup L_7 + e_4 + e_6 & C^{22} &= L_6 \cup L_8 + e_4 + e_5 + e_6 + e_7 \\
 C^{23} &= L_7 \cup L_8 + e_5 + e_7 & C^{24} &= L_1 \cup L_4 + e_1 + e_2 + e_3 \\
 C^{25} &= L_1 \cup L_6 + e_1 + e_5 + e_4 & C^{26} &= L_1 \cup L_7 + e_1 + e_5 + e_6 \\
 C^{27} &= L_2 \cup L_4 + e_3 + e_2 & C^{28} &= L_2 \cup L_6 + e_5 + e_4
 \end{aligned}$$

Their lengths are:

$$\begin{aligned}
 \varepsilon(L_1) &= k - 1, \quad \varepsilon(L_2) = k, & \varepsilon(L_3) &= k, \quad \varepsilon(L_4) = k - 1, \\
 \varepsilon(L_5) &= k - 2, \quad \varepsilon(L_6) = k - 2, & \varepsilon(L_7) &= k - 2, \quad \varepsilon(L_8) = k - 2
 \end{aligned}$$

Thus, we have:

$$\begin{aligned}
 \varepsilon(C^1) &= 2k & \varepsilon(C^{02}) &= 2K + 1 & \varepsilon(C^3) &= 2k & \varepsilon(C^4) &= 2k \\
 \varepsilon(C^5) &= 2k + 1 & \varepsilon(C^6) &= 2k & \varepsilon(C^{07}) &= 2K & \varepsilon(C^{08}) &= 2K
 \end{aligned}$$

$$\varepsilon(C^9) = 2k+1 \quad \varepsilon(C^{010}) = 2K+1 \quad \varepsilon(C^{011}) = 2K+1 \quad \varepsilon(C^{012}) = 2K+1$$

$$\varepsilon(C^{13}) = 2k-1 \quad \varepsilon(C^{14}) = 2k+1 \quad \varepsilon(C^{15}) = 2k+1 \quad \varepsilon(C^{16}) = 2k+1$$

$$\varepsilon(C^{17}) = 2k \quad \varepsilon(C^{018}) = 2K-2 \quad \varepsilon(C^{019}) = 2K \quad \varepsilon(C^{20}) = 2k$$

$$\varepsilon(C^{21}) = 2k-2 \quad \varepsilon(C^{22}) = 2k \quad \varepsilon(C^{23}) = 2k-2 \quad \varepsilon(C^{024}) = 2K+1$$

$$\varepsilon(C^{25}) = 2k \quad \varepsilon(C^{26}) = 2k \quad \varepsilon(C^{027}) = 2K+1 \quad \varepsilon(C^{028}) = 2K$$

It is clear that

$$\varepsilon(C^i) = 2k, i = 1, 3, 4, 6, 7, 8, 17, 19, 20, 22, 25, 26, 28,$$

$$\varepsilon(C^i) = 2k+1, i = 2, 5, 9, 10, 11, 12, 14, 15, 16, 24, 27.$$

$\varepsilon(C^i) = 2k-1, i = 13, 18, 21, 23$ . Now, for any pair  $w$  and  $z$  of  $V(G)$ , are include at least one cycle of the cycles defined above So, we have

$$\text{Max } \{q(C^i) : 1 \leq i \leq 28\} \leq \left\lfloor \frac{2k+1}{2} \right\rfloor = k$$

We get  $p(7, 8k-10) \leq q(G) \leq k$ , for any positive integers  $K$ . Now, From

Deng. [5] we have  $p(7, q) \geq \left\lceil \frac{q}{t+1} \right\rceil = k$ , for  $e = 7$  and  $q = 8k-10$ .

$$p(7, 8k-10) \geq \left\lceil \frac{8k-10}{8} \right\rceil = k \text{ Thus, } p(7, 8k-10) = k.$$

**Proposition 2.2 :**  $p(7, 8k-6) = k$  for  $k \geq 3$

*Proof :* From theorem 2.1 above we have  $p(7, 8k-10) = k$  for  $k \geq 3$ . From Schoon.[7]  $p(e, q') \geq p(e, q)$  if  $q' \geq q$  we get  $p(7, q') \geq p(7, q)$  Where  $q' = 8k-6$

and  $k \geq 3$ , Thus,  $p(7, q') \geq k$ . From A.A.Najim. [9]  $\left\lceil \frac{q-2}{e+1} \right\rceil \leq p(e, q) \leq \left\lfloor \frac{q-2}{e+1} \right\rfloor + 1$

for  $e = 7$  and  $q' = 8k-6$ . So,  $p(e, q') \leq \left\lceil \frac{q'-2}{e+1} \right\rceil + 1 = \left\lceil \frac{8k-6-2}{8} \right\rceil + 1$ . We have

$$p(7, q') \leq k. \text{ Thus } p(7, q') = k$$

**Theorem 2.3 :**  $f(4, q) = 5q-4$  where  $q \geq 4$

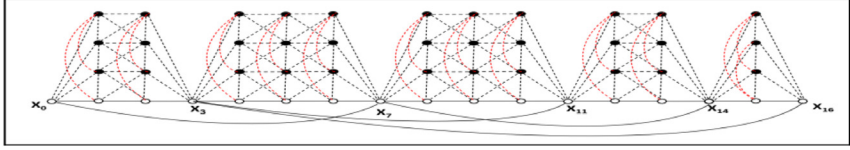


Fig (2.3)

Constriction of  $J_{4,q}$  for  $q \geq 4$ .

*Proof :* We first prove  $f(4,q) \leq 5q - 4$ , let  $G$  be any graph with diameter  $q$ , and let  $G'$  be a connected graph obtained by deleting four edges from  $G$ . From AL-Bachary [11] ( we have  $p(4,5k - 4) = k$  for  $k \geq 2$ ). Let  $q'$  the diameter of  $G'$ ;  $d(G') = q'$  and let  $x$  and  $y$  be two vertices of  $G'$  such that  $d_{G'}(w, z) = q'$  and let be a shortest path  $(x = x_0, x_1, \dots, x_{q'} = y)$  of length  $q'$  in  $G'$ . Partition the vertices in  $V(G')$  in tosets  $V_i$ ,  $i=0,1,2,\dots,q'$  ( $0 \leq i \leq q'$ ) by setting  $V_i = \{u \mid d(x,u) = i\}$ . Notice that all these sets are non empty. Let  $J$  and  $J'$  be the graphs obtained from  $G$  and  $G'$  respectively by contracting eachset  $V_i$  to a single vertex and removing any loops and duplicate edges. Let  $d(J) = j$  and  $d(J') = j'$ . Then  $j \leq q$  and  $j' \leq q'$ . Since  $J'$  simply consist of the path  $x_0, x_1, \dots, x_{q'}$ , we have  $J' = q'$ . The graph  $J$  contains the path  $x_0, x_1, \dots, x_{q'}$ , and at most four additional edges(see figure 2.3 for  $q = 4$ ). Form Wu [3] we have  $j = p(t, f(t, q))$ , then  $j = p(4, q') = \left\lceil \frac{q'+4}{5} \right\rceil \geq \frac{q'+4}{5}$ . So,  $\frac{q'+4}{5} \leq j$ , then  $q' \leq 5j - 4$  where  $(j \leq q)$ .  $q' \leq 5j - 4$  Thus  $f(4,q) \leq 5q - 4$ .

Now, wanted to prove  $f(4,q) \geq 5q - 4$ , from schoone [7] we have  $f(e,q) \geq (e+1)q - e = 5q - 4$ , if  $q$  is even and From Wu and Xu [10]  $f(e,q) = (e+1)q - 2e + 4 = 5q - 4$ , if  $q$  is odd .

Thus  $f(4,q) \geq 5q - 4$ .  $f(4,q) = 5q - 4$

**Theorem 2.4 :**  $f(7,q) = 8q - 10$  where  $q \geq 5$

*Proof :* We want to prove that  $f(7,q) \geq 8q - 10$ , we only need to construct a 8-edge-connected graph with diameter  $q$  such that its diameter increases to at least  $8q - 10$  when its seven edge are deleted (see[12]) . Let  $R_{7,q}$  be a graph obtaining from a path of length  $(8q - 10)$ ,  $p = x_0, x_1, \dots, x_{8q-10}$ , plus seven extra edges:  $e_1 = x_0 x_{2q-1}$ ,  $e_2 = x_{q-1} x_{3q-1}$ ,  $e_3 = x_{2q-1} x_{4q-2}$ ,  $e_4 = x_{q-1} x_{5q-4}$

$e_5 = x_{2q-1}x_{6q-6}$ ,  $e_6 = x_{q-1}x_{7q-8}$ ,  $e_7 = x_{2q-1}x_{8k-10}$ . The prove is similarity to the prove of theorem 2.3

**Theorem 2.5 :**  $T_F(f, q) = \left\lfloor \frac{f-q-3}{q+1} \right\rfloor$

*Proof :* From Chung and Garey in 1984 [4] we have  $F(e, q) \leq q(e+1) - e + 3$

Put  $f = F(e, q)$ , So,  $f \leq q(e+1) - e + 3$ .  $f - q - 3 \leq (q+1)e$ , then  $t \geq \frac{f-q-3}{q-1}$ , Thus  $T_F \geq \left\lfloor \frac{f-q-3}{q-1} \right\rfloor$

**Theorem 2.6 :**

$$\left\lfloor \frac{f-6}{q-2} \right\rfloor - 1 \geq T_F(f, q) \geq \begin{cases} \left\lfloor \frac{f-q}{q+1} \right\rfloor & \text{iff } (e, q) \leq (e+1)q + e \\ \left\lfloor \frac{f-2}{q+2} \right\rfloor - 1 & \text{iff } (e, q) \leq (e+1)q - e + 1 \end{cases}$$

*Proof :* from Wu and Xu in 2006 [10] we have  $f(e, q) \geq (e+1)q - 2e + 4$  for  $e \geq 4$  and  $q \geq 3$  is odd

Put  $f = f(e, q)$ , So  $f \geq (e+1)q - 2e + 4$

$$f - q - 4 \geq e(q - 2)$$

$$e \leq \frac{f-q-4}{q-2} \text{ then } e \leq \frac{f-6}{q-2} - 1 \quad e \leq \frac{f-6}{q-2} - 1$$

Thus  $T_F \leq \left\lfloor \frac{f-6}{q-2} \right\rfloor - 1$  From A.A. Najim and J.M. Xu 2006 [8] we have

$$f(e, q) \leq \begin{cases} (e+1)q & \text{if } p(e, q) = \frac{q-2}{e+1} \\ (e+1)q - e + 1 & \text{if } p(e, q) = \frac{q-2}{e+1} + 1 \end{cases}$$

for  $e \geq 4$  and  $q \geq 3$ .

Now, if  $f(e, q) \leq (e+1)q + e$

Put  $f = f(e, q)$  we get  $f \leq (e+1)q + e$ , then  $e \geq \frac{f-q}{q+1}$ , Thus  $T_F \geq \left\lfloor \frac{f-q}{q+1} \right\rfloor$

If  $f(e, q) \leq (e+1)q - e + 1$

Put  $f = f(e, q)$  we get  $f \leq (e + 1)q - e + 1$

$f - q - 1 \leq e(q - 1)$ , So,  $e \geq \frac{f - q - 1}{q - 1}$  Thus  $T_F \geq \left\lfloor \frac{f - 2}{q - 1} \right\rfloor - 1$

So

$$\left\lfloor \frac{f - 6}{q - 2} \right\rfloor - 1 \geq T_F(f, q) \geq \begin{cases} \left\lfloor \frac{f - q}{q + 1} \right\rfloor & \text{iff } (e, q) \leq (e + 1)q + e \\ \left\lfloor \frac{f - 2}{q + 2} \right\rfloor - 1 & \text{iff } (e, q) \leq (e + 1)q - e + 1 \end{cases}$$

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